



Chapter 1. Introduction

System Programming

<http://www.cs.ccu.edu.tw/~pahsiung/courses/sp>

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Introduction

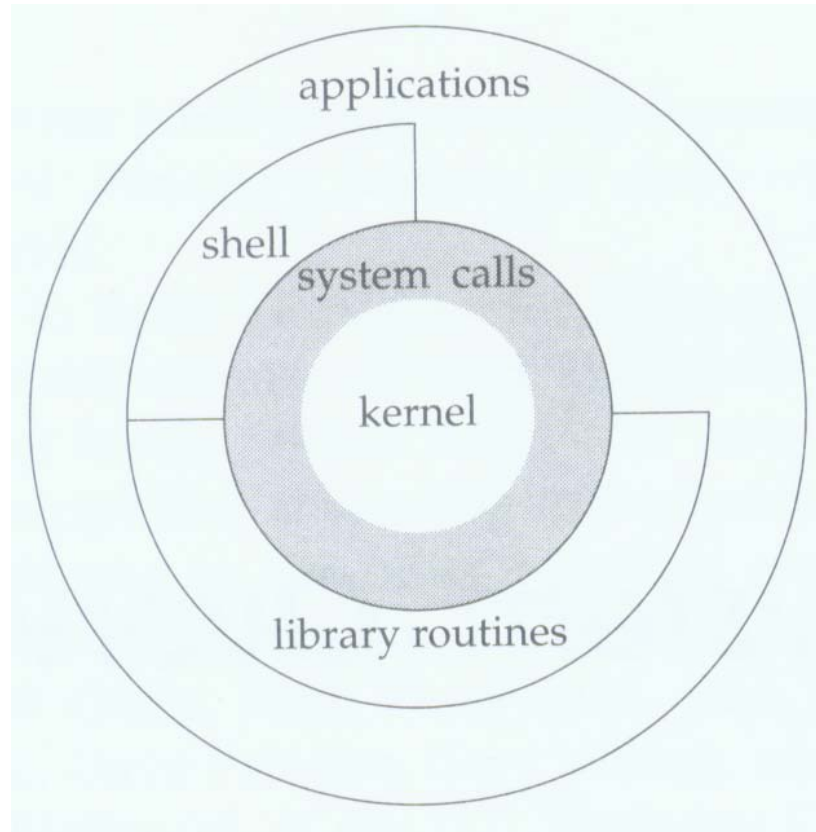
- Whirlwind tour of UNIX from a programmer's perspective
- Brief descriptions and examples
- Services provided by UNIX



UNIX Architecture

- Operating system is the software that
 - controls hardware resources
 - Provides program execution environment
- Architecture
 - Kernel
 - System Calls
 - Shell
 - Applications

Architecture of UNIX OS





Logging In

- Login name
- Password
- Colon-separated fields in each entry of `/etc/passwd`
 - name,
 - encrypted password or “x” (password is in `/etc/shadow`),
 - numeric user ID,
 - numeric group ID,
 - real name,
 - home directory,
 - shell program
- `sar:x:205:105:Stephen Rago:/home/sar:/bin/ksh`



Logging In

- Shells: a command line interpreter that reads user input and executes commands
 - Bourne shell: `/bin/sh` (used by root)
 - C shell: `/bin/csh` (used by users)
 - Korn shell: `/bin/ksh` (successor of Bourne shell)
 - Bourne-again shell: `/bin/bash` (all Linux systems)
 - TENEX C shell: `/bin/tcsh` (replacement of C shell)



Files and Directories

- Filesystem: hierarchical arrangement of directories and files
- Root directory: /
- File attributes: type, size, owner, permissions, last modification time, ...
- `stat()`, `fstat()`: return file attribute struct



Files and Directories

- Filename
 - Chars not allowed: (/) and (NULL)
- Two filenames automatically created whenever a new dir is created:
 - . → current directory
 - .. → parent directory
- What is .. in root directory (/)?



Files and Directories

- Pathname
 - A sequence of zero or more filenames, separated by slashes (/), and optionally starting with a slash
- Absolute pathname
- Relative pathname



Program 1.3:

(bare bones implementation of ls command)

```
#include <sys/types.h>
#include <dirent.h>
#include "apue.h"

int
main(int argc, char *argv[])
{
    DIR          *dp;
    struct dirent *dirp;

    if (argc != 2)
        err_quit("a single argument (the directory name) is required");

    if ( (dp = opendir(argv[1])) == NULL)
        err_sys("can't open %s", argv[1]);

    while ( (dirp = readdir(dp)) != NULL)
        printf("%s\n", dirp->d_name);

    closedir(dp);
    exit(0);
}
```



Program 1.3

- Edit and save in myls.c

```
cc myls.c      (output: a.out)
```

```
./a.out /dev (output: ... ..)
```

```
./a.out /var/spool/mqueue
```

```
can't open /var/spool/mqueue:  
Permission denied
```

```
./a.out /dev/tty
```

```
can't open /dev/tty: Not a directory
```



Program 1.3 (details)

- `apue.h` (in Appendix B)
- ANSI C function declaration
- `argv[1]`: is an argument to our prog
- What is `argv[0]`?
- `opendir()`, `readdir()`, `closedir()`
- `DIR`: directory struct
- `err_sys()`, `err_quit()`: error routines



Program 1.3 (details)

- `exit(0)`: program done!
- return value:
 - 0 → OK!
 - 1 ~ 255 → ERROR!
- Current working directory (CWD)
- `chdir()`: change working directory
- Eg: `doc/my/file`: relative to CWD
- Home directory: in password file



Input and Output

- File Descriptors
 - small non-negative integers that kernel uses to identify files being accessed by a process
- Standard Input
- Standard Output
- Standard Error



Input and Output

- ls
 - stdin, stdout, stderr: → terminal
- ls > myfile.abc
 - stdout: myfile.abc
- How to redirect **stderr** to a file???
- How to redirect **stdin** from a file?
- Unbuffered I/O
 - open(), read(), write(), lseek(), close()



Program 1.4: stdin → stdout

```
#include          "apue.h"

#define BUFSIZE   8192

int
main(void)
{
    int          n;
    char buf[BUFSIZE];

    while ( (n = read(STDIN_FILENO, buf, BUFSIZE)) > 0)
        if (write(STDOUT_FILENO, buf, n) != n)
            err_sys("write error");

    if (n < 0)
        err_sys("read error");

    exit(0);
}
```




Program 1.4 (details)

- unistd.h: constants are defined:
 - STDIN_FILENO = 0
 - STDOUT_FILENO = 1
- BUFSIZE constant:
 - affects program efficiency
- read() returns:
 - #bytes read
 - 0 if EOF
 - -1 if error



Standard I/O

- A buffered interface
- No need to worry about BUFSIZE
- Deal with “lines of input”
 - fgets() reads an entire line
 - read() reads a specified # of bytes
- printf() (#include <stdio.h>)



Program 1.5: stdin → stdout using standard I/O

```
#include      "apue.h"

int
main(void)
{
    int      c;

    while ( (c = getc(stdin)) != EOF)
        if (putc(c, stdout) == EOF)
            err_sys("output error");

    if (ferror(stdin))
        err_sys("input error");

    exit(0);
}
```



Program 1.5 (details)

- `getc()` reads 1 char at a time
- `putc()` writes 1 char at a time
- `#include <stdio.h>`
 - `stdin`: standard input
 - `stdout`: standard output



Programs and Processes

- Program: an **executable file** on disk
- Process: an **executing instance** of a program
- Process also called “task” by some OS
- Unique non-negative integer identifier for each process (pid)



Program 1.6: process ID

```
#include "apue.h"

int
main(void)
{
    printf("hello world from process
    ID %d\n", getpid());
    exit(0);
}
```



Process Control

- Three functions
 - `fork()`
 - `exec()`: 6 variants
 - `waitpid()`



Program 1.7: exec stdin cmds

```
#include <sys/types.h>
#include <sys/wait.h>
#include "apue.h"
int
main(void)
{
    char  buf[MAXLINE];
    pid_t pid;
    int   status;

    printf("%% "); /* print prompt (printf requires %% to print %) */
    while (fgets(buf, MAXLINE, stdin) != NULL) {
        buf[strlen(buf) - 1] = 0; /* replace newline with null */
        if ( (pid = fork()) < 0)
            err_sys("fork error");
        else if (pid == 0) { /* child */
            execlp(buf, buf, (char *) 0);
            err_ret("couldn't execute: %s", buf);
            exit(127);
        }
        /* parent */
        if ( (pid = waitpid(pid, &status, 0)) < 0)
            err_sys("waitpid error");
        printf("%% ");
    }
    exit(0);
}
```




Program 1.7 (details)

- `fgets()` reads a newline-terminated line
- replace newline with NULL using `strlen()`
- `fork()` creates a new process
- `fork()` returns pid of new child process to parent
- `fork()` returns 0 to child
- `execlp()` in child executes stdin cmd



Program 1.7 (details)

- parent waits for child to finish
- `waitpid()`
- current version:
can't pass arguments to the stdin cmd



Threads

- A process can have one or more threads
 - Can exploit parallelism on multiprocessor systems
 - Same address space, file descriptors, stacks, process-related attributes
 - Need to synchronize access to shared data
 - Thread IDs: local to process



Error Handling

- Negative return value when error occurs
- `#include <errno.h>`
- `errno` variable
 - never cleared if error does not occur
 - never set to 0 by any function



Error Handling (contd)

- 2 functions for printing error messages:

```
#include <string.h>  
char *strerror(int errnum);
```

```
#include <stdio.h>  
void perror(const char *msg);
```

- strerror() returns a string
- perror() outputs “msg: <error_msg>”



Program 1.8: use of error func

```
#include    <errno.h>
#include    "apue.h"

int
main(int argc, char *argv[])
{
    fprintf(stderr, "EACCES: %s\n",
        strerror(EACCES));

    errno = ENOENT;
    perror(argv[0]);

    exit(0);
}
```



Program 1.8: results

```
$ a.out
```

```
EACCES: Permission denied
```

```
a.out: No such file or directory
```

- `prog1 < inputfile | prog2 | prog3 > outputfile`
- `(argv[0] passed as arg to perror())`



Error Recover

- Fatal error: no recovery action
- Nonfatal error: delay and try again
 - Improves robustness by avoiding an abnormal exit
 - Examples
 - EAGAIN, ENFILE, ENOBUFS, ENOLCK, ENOSPC, ENOSR, EWOULDBLOCK, ENOMEM



User Identification

- User ID: numeric identifier of a user
- Group ID: numeric identifier of a group

```
#include "apue.h"

int
main(void)
{
    printf("uid = %d, gid = %d\n",
           getuid(), getgid());
    exit(0);
}
```



Signals

- A technique to notify a process that some condition has occurred
- E.g.: divide by zero → SIGFPE
- Process response to a signal
 - Ignore the signal, OR
 - Let the default action occur, OR
 - Provide a function to handle the signal.



Program 1.10

```
#include "apue.h"
#include <sys/wait.h>

static void sig_int(int); /* our signal-catching function */

int
main(void)
{
    char buf[MAXLINE];
    pid_t pid;
    int status;

    if (signal(SIGINT, sig_int) == SIG_ERR)
        err_sys("signal error");

    printf("%s "); /* print prompt (printf requires %s to print %) */
    while (fgets(buf, MAXLINE, stdin) != NULL) {
        buf[strlen(buf) - 1] = 0; /* replace newline with null */

        if ( (pid = fork()) < 0 )
```



Program 1.10

```
        err_sys("fork error");

    else if (pid == 0) {                /* child */
        execlp(buf, buf, (char *) 0);
        err_ret("couldn't execute: %s", buf);
        exit(127);
    }

    /* parent */
    if ( (pid = waitpid(pid, &status, 0)) < 0)
        err_sys("waitpid error");
    printf("%% ");
}
exit(0);
}

void
sig_int(int signo)
{
    printf("interrupt\n%% ");
}
}
```



Time Values

- Two different time values
- **Calendar time:** #seconds since the Epoch, which is 00:00:00 Jan 1, 1970, Coordinated Universal Time (UTC).
- **Process time:** measures CPU resources used by a process, in clock ticks, which is 50, 60, or 100 ticks per second.



UNIX Time Values (contd)

- Execution time of a process has 3 values:
- **clock time**: total amount of time from process start to finish
- **user CPU time**: CPU time due to user instructions in a process
- **system CPU time**: CPU time due to kernel activities on behalf of the process



UNIX Time Values (contd)

- To measure process execution time, use the “time” command as follows:

```
$ time ls > /dev/null
```

```
real    0m19.81s
```

```
user    0m0.43s
```

```
sys     0m4.53s
```



System Calls & Library Functions

- **System Calls:**
 - Entry points into an OS kernel
- Cannot be changed by user
- A function of the same name in the standard C library
- User just calls those C functions whenever system calls are needed



System Calls & Library Functions

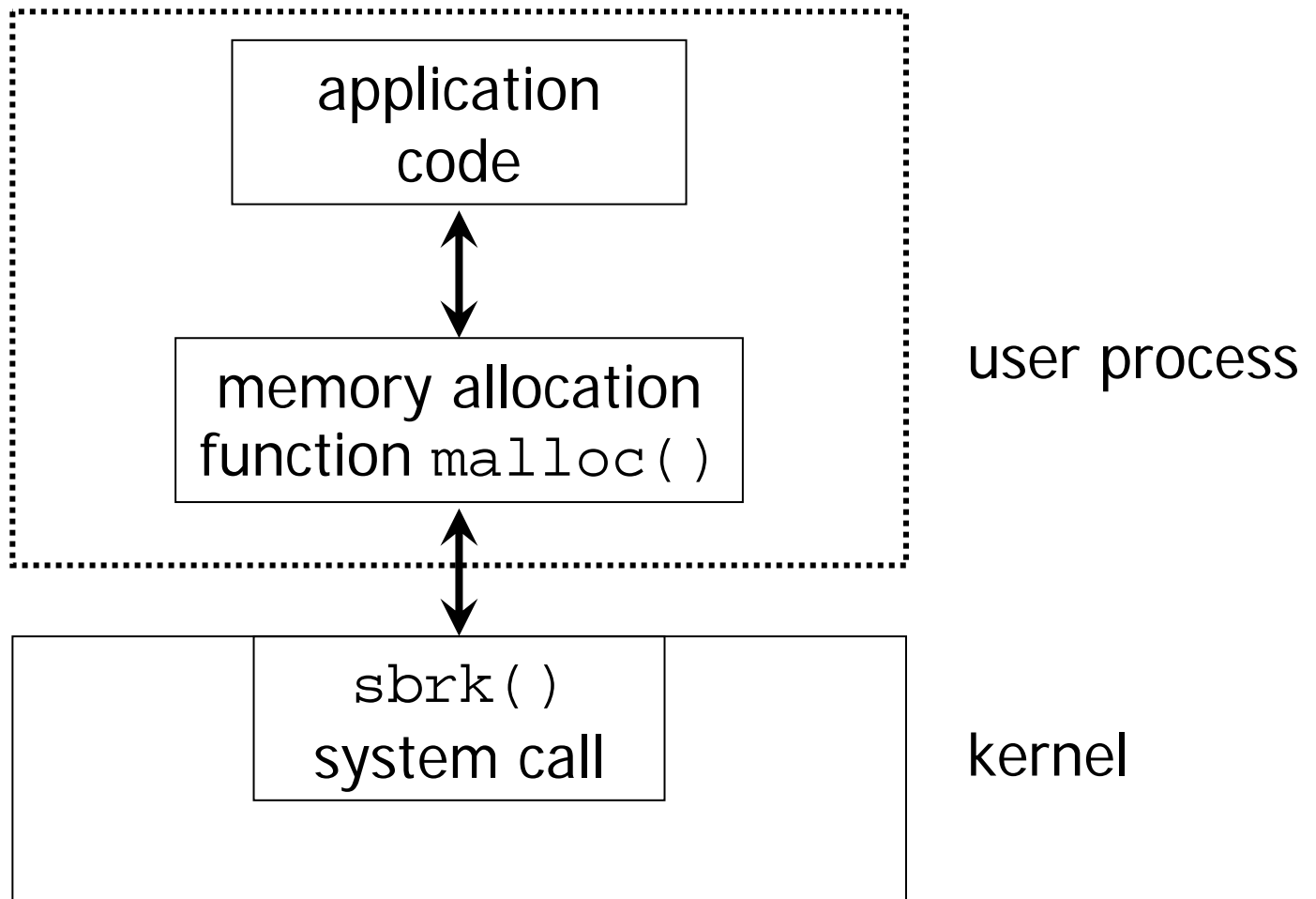
- Library Functions: not entry points into kernel, just functions, but they may invoke one or more system calls
- E.g.: `printf()` invokes `write()` system call
- E.g.: `strcpy()`, `atoi()`: do not invoke any system call
- Implementor view: fundamental diff
- Programmer view: no critical difference



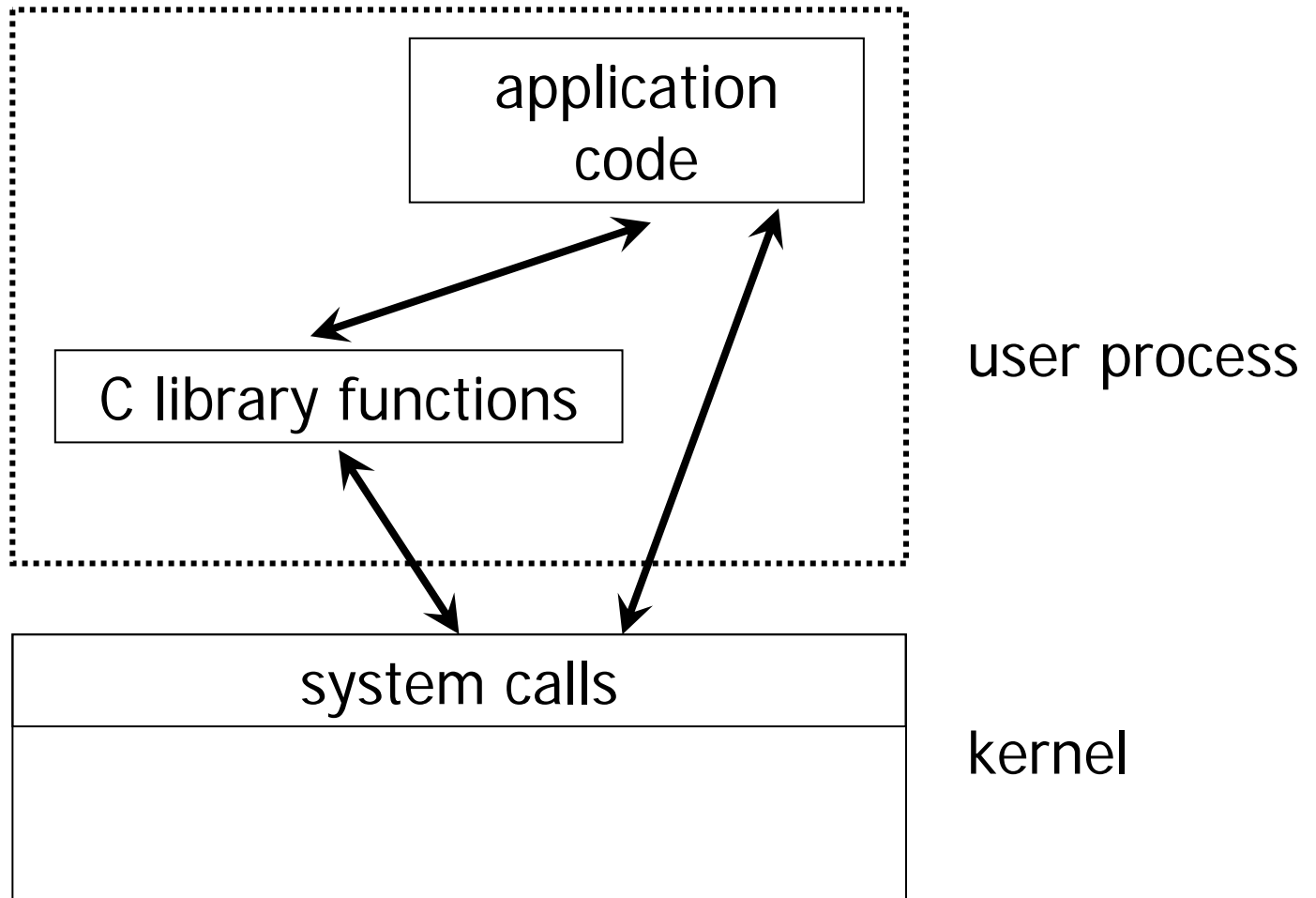
System Calls & Library Functions

- `malloc()`: Memory allocation
- Many ways to do memory allocation and garbage collection (best fit, first fit)
- `sbrk()`: UNIX system call, increases or decreases address space of process by a specified number of bytes
- Can implement own mem alloc function using `sbrk()`

System Calls & Library Functions



System Calls & Library Functions





System Calls & Library Functions

- **System call**: time in #seconds since Jan 1, 1970
- **C function**: human-readable time and date using local time zone
- Both are called “**functions**” in the textbook