

## II.4. Feedback of Rate and Loss Information for Networked Video

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### Outline

- Motivation for feedback and adaptation
- Transport of video over the best-effort internet
- Source adaptation using rate-based feedback control
- Rate characterization and smoothing
- The SAVE algorithm
- Evaluation criteria
- Experiments and simulation results
- Conclusion

### Motivation for Feedback and Adaptation

- Three drawbacks of an open-loop transport of video :
  - The requirement for conservative admission control
  - Large buffers (impacting delay)
  - The potential for packet loss (impacting quality)

### Motivation for Feedback and Adaptation (Cont.)

### Transport of Video over the Best-Effort Internet

- Two types of transporting video over the best-effort internet :
  - Unicast (Point to Point)
  - Multicast (Point to Multipoint)

### The AIMD Rate Control Algorithm

If (Congested)  
 $\lambda = \max\{\lambda \times \alpha, \text{minimum\_rate}\};$   
If (UnCongested)  
 $\lambda = \max\{\lambda + \beta, \text{minimum\_rate}\};$   
  
 $\lambda$ : source rate;  
 $\alpha$ : reduction factor ( $\leq 1$ );  
 $\beta$ : increase factor

### Transport of Video over the Best-Effort Internet

- RAP : Involves a rate-adaptive source at a server and a corresponding rate-adaptive receiver at every client, with the receiver acknowledging every packet
- The interpacket gap (IPG) is adjusted in RAP in response to congestion. The adjusting strategy is the additive increase, multiplicative decrease
- Real-time streams that do not use TCP (e.g., they may use UDP) may gain an unfair advantage over TCP-transported data
  - Possible Solution: "TCP-friendly" protocols

### Supporting Video in A Multicast Environment

- McCanne proposed to move the rate adaptation to the receivers by means of receiver-driven layered multicast (RLM).
  - Combines layered video compression with a layered transmission scheme
  - RLM receivers use join experiments to add a new layer at well chosen times
  - RLM groups use shared learning to improve join experiments.

### Source Adaptation Using Rate-Based Control

- Constraints on rate required for transporting real-time video
- Rate renegotiation mechanisms and adaptation

### Rate Constraints for Transporting Real-time Video

- Buffering and delay constraints: (Buffer underflow or overflow)
  - End-to-end delay
- Source adaptation constraints
  - Received quality, delays, ...
- Lookahead constraints (e.g., for the stored video)
- Implicit versus explicit feedback
- Signaling frequency and latency constraints
- Rate prediction error
  - Uncertainties on the network and source rates

### Rate Renegotiation Mechanism and Adaptation

- RCBR uses renegotiation of the traffic parameters of a constant bit rate connection by means of ATM signaling messages.
- The control mechanism proposed by Kanakia et al. (3b) is based on predicting the evolution of the system over time and using that prediction to compute a target sending rate for each frame of video data.

### Rate Renegotiation Mechanism and Adaptation

If  $X_{n,k} = 0$   
 $\lambda_n = \lambda_{n-1} + \delta$   
 else  
 $\lambda_n = \mu_n + (X^* - X_n) / (\text{gain} * F)$   
 $X^*$  : The target value of the buffer occupancy at the bottleneck for this video flow.  
 $1/F$ : The frame rate for this video  
 $\mu_n$  : The service rate  
 $X_n$ : The queue occupancy

## Rate Characterization and Smoothing

$S(t) = \sum s_i$  is the cumulative amount of data sent up to time  $t$

$l_i$  denote the sizes of frame in a stream of  $N$  video frames

To avoid of underflow requires that  $L(t) \leq S(t)$

To avoid of overflow requires that

$$S(t) \leq U(t) := \min\{L(t-1) + B, L(N)\}$$

## The SAVE Algorithm

- SAVE (Smoothed Adaptive Video over Explicit rate networks) is a source algorithm used for transporting compressed video in conjunction with explicit rate-based feedback control in the network.
- This algorithm comprises two parts
  - The rate request algorithm : specifies how the source requests bandwidth from the system
  - The frame quantization algorithm : specifies how the frame size are controlled to avoid excessive delay

## Heuristics for the Requested Rate

- Because of the GOP structure, one determinant of the required rate will be the smoothed rate :

$$r_{sm} = f_{sm} / \tau$$

$\tau$  is the interframe time

$f_{sm}$  is the ideal frame size that required by the encoder to encode the frame at ideal perceptual lossless quality

## Rate Request & Frame Quantization Algorithms

$$r_{sm}(n) = (\tau w_{sm})^{-1} \sum_{i=0}^{w_{sm}-1} f(n-i)$$

$$r_{max}(n) = (\tau w_{max})^{-1} \max f(n-i)$$

$$i = 0, 1, \dots, w_{max}-1$$

➤ The requested rate at frame  $n$  is :

$$r_{req}(n) = \beta \max \{r_{sm}(n), r_{max}(n), r_{ar}(n)\}$$

➤ The encoded size of frame  $n$  is :

$$f_{enc}(n) = \min \{f(n), \max\{f_{avail}(n-1), r \times f(n)\}\}$$

## Evaluation Criteria

- Source delay
- Quality and Adaptation
- Robustness to Network Feedback Delay
- Channel Capacity and Multiplexing Gain
- Sensitivity to Algorithm Parameters

## Evaluation Criteria (Cont.)

- Source delay :  
The primary concern for the work is the aggregate delay introduced in the source buffer and the network.
- Quality and Adaptation :  
Cropping : The reduction from the ideal rate to the encoded rate.  
Evaluate the pattern of cropping over the entire sequence of frames.

## Evaluation Criteria (Cont.)

- **Robustness to Network Feedback Delay :**  
The rate allocation mechanism of the explicit rate network is not expected to instantaneously allocate a rate in response to requests.
- **Channel Capacity and Multiplexing Gain :**
- **Sensitivity to Algorithm Parameters**

## Experiments and Simulation Results

- Modeling network characteristics with frame-level simulations  
if  $n > \delta$   
$$r_{\text{all}}(n) = r_{\text{req}}(n - \delta)$$
  
else  
$$r_{\text{all}}(n) = r_0$$

## Experiments and Simulation Results

- Modeling network characteristics with frame-level simulations  
$$p(n) = \max \{1, C/R(n)\}$$
  
if  $n > \delta$   
$$r_{\text{all}}(n) = P(n - \delta) r_{\text{req}}(n - \delta)$$
  
else  
$$r_{\text{all}}(n) = r_0$$

## Experiments and Simulation Results

- Modeling network characteristics with frame-level simulations  
The buffer evolution is :  
$$b(n) = f_{\text{enc}}(n) + \max\{0, b(n-1) - r_{\text{all}}(n-1) \times \tau\}$$

## Network Configuration of Cell-Level Simulation

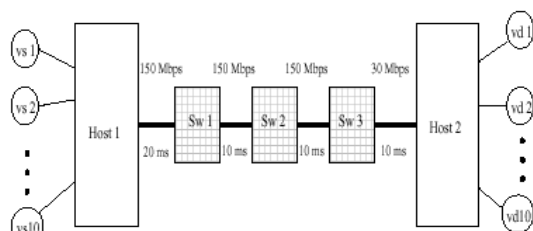
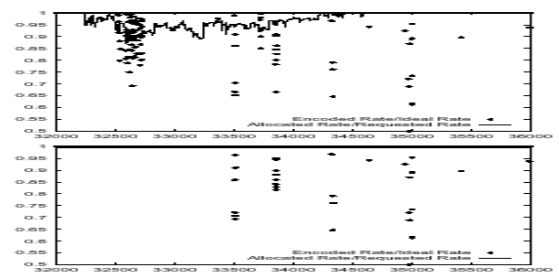
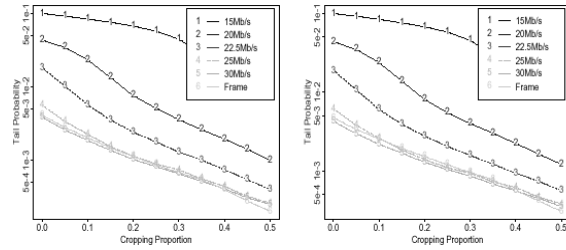


Figure 3: NETWORK CONFIGURATION FOR CELL-LEVEL SIMULATION

## Impact of Network Congestion & Feedback Delay on Quality



## Behavior of Frame Delay



## Behavior of Frame Delay (Cont.)

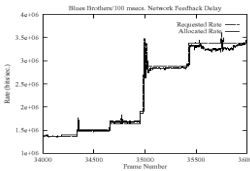


Figure 6: NETWORK SIMULATION RATES. Trace A, 1 trace within a multiplexed set. For short periods, allocated rate is lower than requested due to network contention.

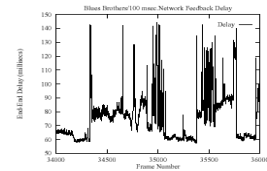


Figure 7: END-TO-END NETWORK DELAY. Trace A, 1 trace within a multiplexed set.

## Multiplexing Gain & Impact on Quality

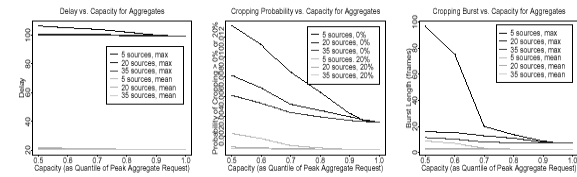


Figure 8: VARIATION WITH AGGREGATION SIZE OF IMPACT OF RATE-REDUCTION THROUGH REDUCED CAPACITY. For aggregations of 5, 20, 35 sources, averaged impact on single traces within aggregate; capacity expressed as quantile of peak aggregate requested rate. TOP: mean and max. delay. MIDDLE: proportion of cropping > 0% and > 20%. BOTTOM: mean and max. burst length of cropping > 20%.

## Multiplexing Gain & Impact on Quality (Cont.)

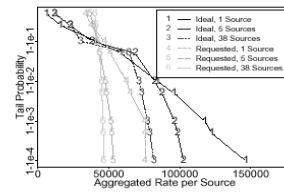


Figure 9: RATE SMOOTHING AND AGGREGATION TAIL distribution, for ideal and requested rates per source, for aggregations of 1, 5 and 38 sources of Set E. High quantiles of requested rate are lower for requested rate.

## Multiplexing Gain & Impact on Quality (Cont.)

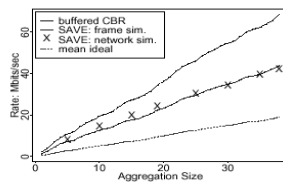


Figure 10: UTILITY OF EXPLICIT RATE FOR MULTIPLEXING GAIN. Bandwidth requirements, normalized by mean ideal rate, for 1- to 38-fold aggregates of Set E. LOWER CURVE: SAVE, 90<sup>th</sup> percentile of aggregate requested rate. UPPER CURVE, CBR rate for 100ms source buffer. Also shown: mean ideal rate, and capacity estimate based on full network simulation described below.

## Multiplexing Gain & Impact on Quality (Cont.)

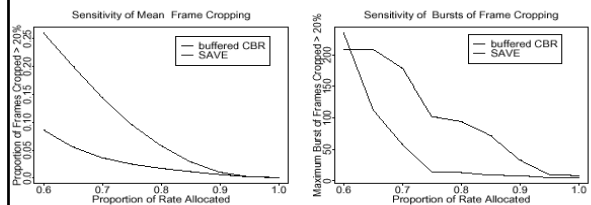


Figure 11: COMPARATIVE SENSITIVITY OF FRAME CROPPING IN SAVE AND BUFFERED CBR TO SYSTEMATIC RATE-REDUCTION OR UNDERALLOCATION OF BANDWIDTH. LEFT: average cropping more sensitive for SAVE. RIGHT: Burst cropping is more sensitive for buffered CBR.

## Sensitivity to Smoothing Parameters

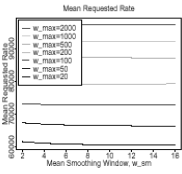


Figure 12: MEAN REQUESTED RATE VS.  $u_{\text{REQ}}$  trace set A, for varying  $u_{\text{REQ}}$ .

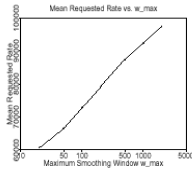


Figure 13: MEAN REQUESTED RATE VS.  $u_{\text{REQ}}$ , averaged over  $u_{\text{REQ}}$ .

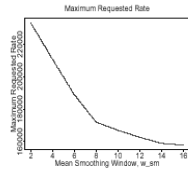


Figure 14: MAXIMUM REQUESTED RATE VS.  $u_{\text{REQ}}$  trace set A. Curves coincide for  $u_{\text{REQ}}$  from 20 to 2000.

## Sensitivity to Smoothing Parameters (Cont.)

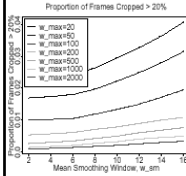


Figure 15: PROPORTION OF FRAMES CROPPED > 20% VS.  $u_{\text{REQ}}$  trace set A for varying  $u_{\text{REQ}}$ .

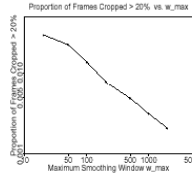


Figure 16: PROPORTION OF FRAMES CROPPED > 20% VS.  $u_{\text{REQ}}$ , trace set A, averaged over  $u_{\text{REQ}}$ .

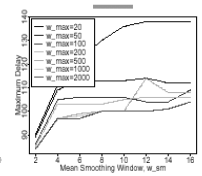


Figure 17: MAXIMUM DELAY VS.  $u_{\text{REQ}}$  trace set A for varying  $u_{\text{REQ}}$ .

## Summary

- Source adaptation using rate-based feedback control
- Rate smoothing
- The SAVE algorithm
- Evaluation criteria